

Rare Event Simulation for the Stationary Distribution of a Markov Chain

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PART I: Multilevel Splitting

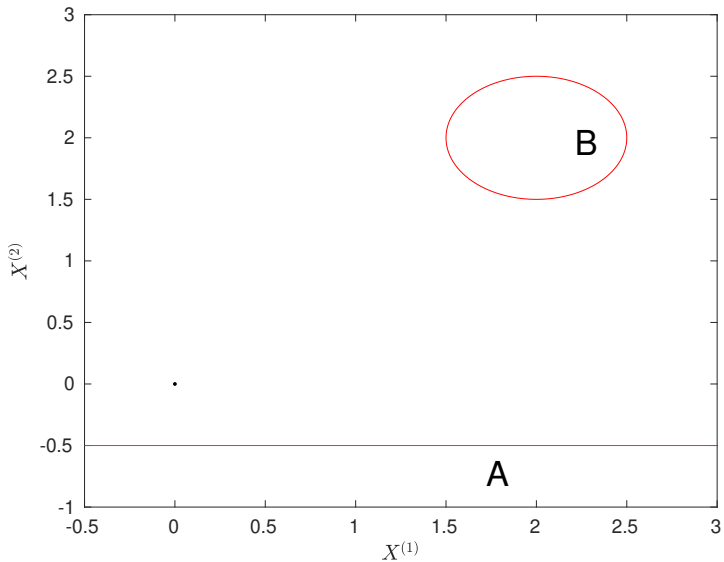
Multilevel Splitting setting

Let $(X_t)_{t \in [0, \infty)}$ be an \mathbb{R}^d -valued Markov process with initial condition $X_0 = \mathbf{0}$. We are interested in finding:

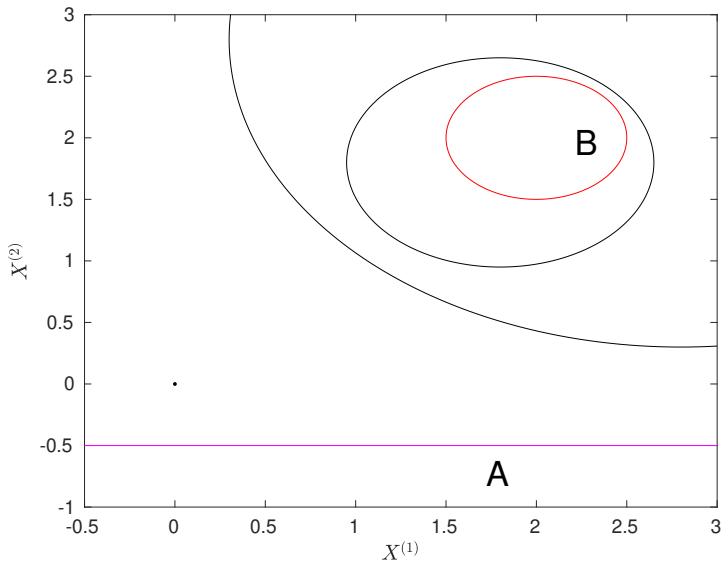
$$\rho := \mathbb{P}(\tau_B < \tau_A),$$

where $\tau_A := \inf\{t > 0 : X_t \in A\}$ and $\tau_B := \inf\{t > 0 : X_t \in B\}$. Here, B is very small and A is some large absorbing set.

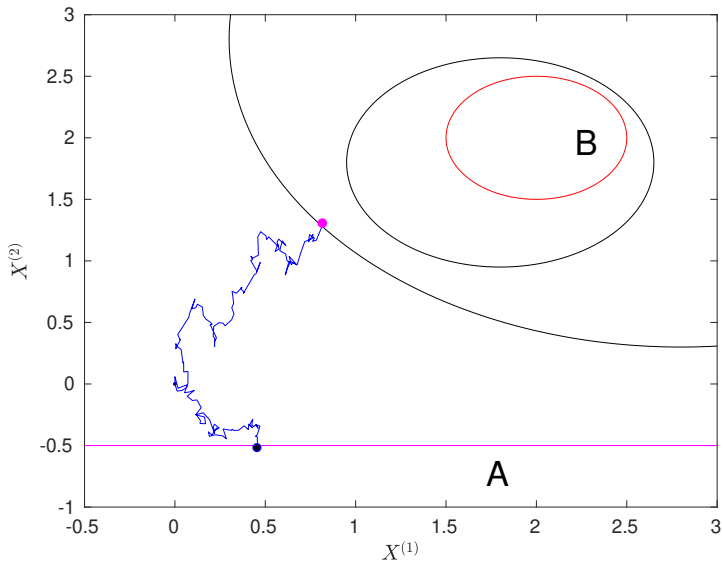
Multilevel Splitting



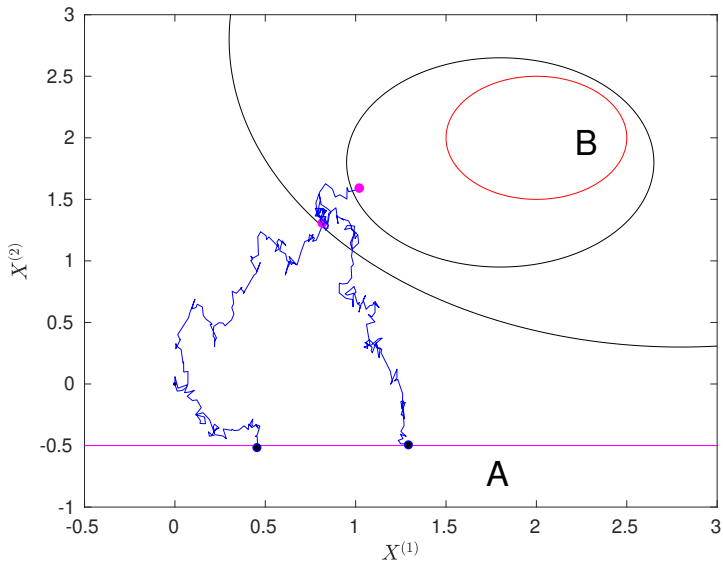
Multilevel Splitting



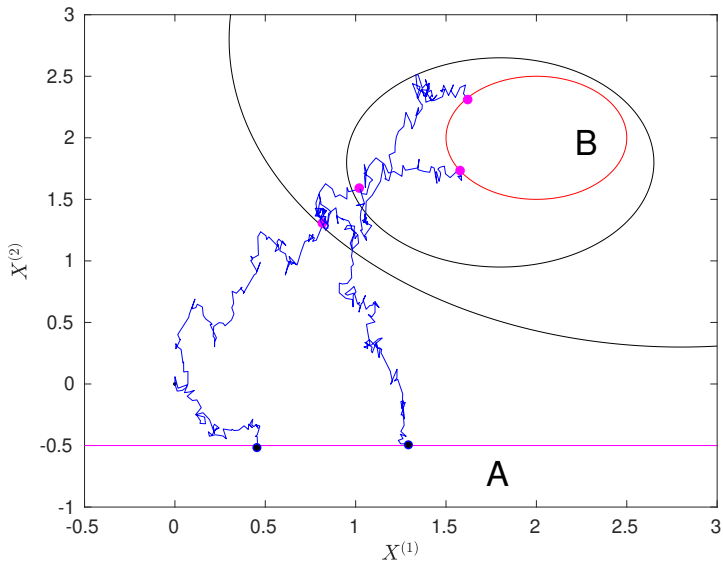
Multilevel Splitting



Multilevel Splitting



Multilevel Splitting



Description of the algorithm

Let n_k be *splitting factors* (in our example $n_k = 2$) and r be the total number of paths that reached set B . We put

$$\hat{p} := \frac{r}{\prod_{k=0}^{m-1} n_k}$$

\hat{p} is an unbiased estimator of p for **ANY** choice of intermediate sets.

For a good choice of intermediate sets, the total computational cost of the estimation is proportional to $(\log p)^2$ instead of p^{-1} !

PART II

Rare Event Simulation for the Stationary Distribution of a Markov Chain

Stationary Markov Chains

Let $X = (X_n)_{n \in \mathbb{N}}$ be an \mathbb{R}^d -valued, time-discrete Markov chain with *stationary (invariant)* measure μ , that is, as $n \rightarrow \infty$,

$$X_n \rightsquigarrow X_\infty \sim \mu.$$

Context: numerical solutions to SDEs:

$$dX_t = f(X_t)dt + g(X_t)dW_t$$

and

$$X_{n+1} = X_n + f(X_n)h + g(X_n)\sqrt{h}\Delta W_n$$

We want to estimate $\mu(B)$ when $\mu(B) \ll 1$. From Ergodic Theorem, for any set B :

$$\mu(B) = \lim_{N \rightarrow \infty} \frac{1}{N} \sum_{n=1}^N \mathbb{1}\{X_n \in B\}.$$

Recurrent Structure of a Markov Chain

Consider the following decomposition of a Markov chain:

- ▶ Choose a set $A \subset \mathbb{R}^d$, with $\mu(A) \in (0, 1)$.
- ▶ Let S_k be consecutive *inwards crossings* of set A , with $S_{-1} = 0$,

$$S_k := \inf\{n > S_{k-1} : X_{n-1} \notin A, X_n \in A\}.$$

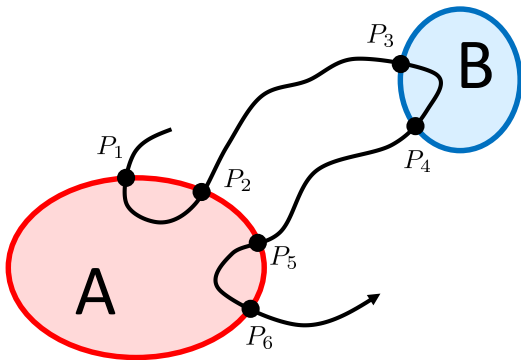
- ▶ Let \mathcal{C}_k be a path within k th cycle, i.e.

$$\mathcal{C}_k := (X_n : S_{k-1} \leq n < S_k).$$

- ▶ We distinguish the cycle *length* and *origin*

$$L_k := S_k - S_{k-1}, \quad X_k^A := X_{S_{k-1}}$$

Note: Assuming the chain is ‘sufficiently nice’, we have $\mathbb{E}L_k < \infty$.



The cycle begins at P_1 and ends at P_5 .

Recurrent Structure of a Markov Chain

In the stationary regime, the cycles $\mathcal{C}_1, \mathcal{C}_2, \dots$ are **identically distributed**.

Define the *time spent in set B* within a cycle:

$$R_k := \sum_{n=S_{k-1}}^{S_k-1} \mathbb{1}\{X_n \in B\}.$$

The quantity of interest $p = \mu(B)$ can be expressed as:

$$p = \frac{\mathbb{E}R_1}{\mathbb{E}L_1}.$$

- ▶ Estimation of $\mathbb{E}L_1$ is easy (Monte Carlo).

Estimation of $\mathbb{E}R_1$

Notice that

$$T_B := \mathbb{E}R_1 = \mathbb{P}(\tau_B < \tau_A^{\text{in}}) \cdot \mathbb{E}(R_1 \mid R_1 > 0)$$

This fits in the framework of MLS with an extra stage with splitting factor n_m !

$$\hat{T}_B := \frac{r_m}{\prod_{k=0}^{m-1} n_k} \cdot \frac{\sum_{j=1}^{n_m r_m} \hat{R}_+^{(j)}}{n_m r_m}$$

\hat{T}_B is an unbiased estimator for $\mathbb{E}R_1$!

For a good choice of intermediate sets and **set A**, the computational cost is proportional to $(\log p)^2$!

Assumptions leading to optimality

The study of efficiency of the algorithm is intractable for a general choice of intermediate sets. Recall that

$$\tau_k = \inf\{n > 0 : X_n \in B_k\}, \quad D_k := \{\tau_k < \tau_A\}$$

We assume the following:

(I) for all $k \in \{1, \dots, m-1\}$, for all X_{τ_k} ,

$$\mathbb{P}(D_{k+1} \mid D_k, X_{\tau_k}) \equiv \mathbb{P}(D_{k+1} \mid D_k)$$

(II) for all X_1^A ,

$$\mathbb{P}(\tau_B < \tau_A \mid X_1^A) \equiv \mathbb{P}(\tau_B < \tau_A)$$

(III) for all cycle origins X_{τ_B} ,

$$(R_1 \mid R_1 > 0, X_{\tau_B}) \stackrel{d}{=} (R_1 \mid R_1 > 0) =: R_+,$$

Numerical Example: Franzke (2012) Model

Franzke (2012) Model

$$dx_1 = \mu \left(-x_2(L_{12} + a_1x_1 + a_2x_2) + d_1x_1 + F_1 \right. \\ \left. + L_{13}y_1 + B_{123}^1x_2y_1 + (B_{131}^2 + B_{113}^2)x_1y_1 \right) dt$$

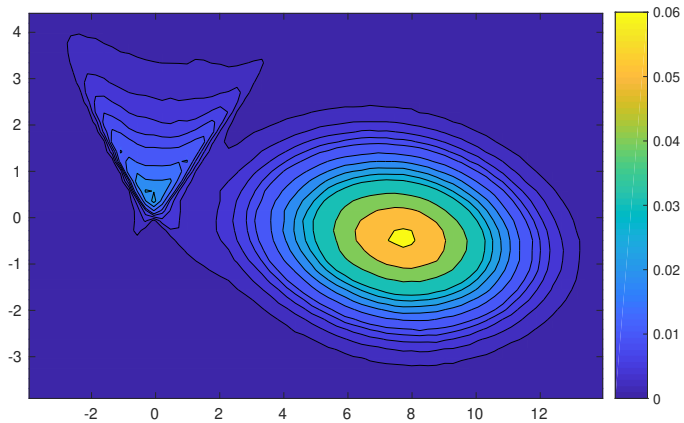
$$dx_2 = \mu \left(+x_1(L_{21} + a_1x_1 + a_2x_2) + d_2x_2 + F_2 \right. \\ \left. + L_{24}y_2 + B_{213}^1x_1y_1 + (B_{242}^3 + B_{224}^3)x_2y_2 \right) dt$$

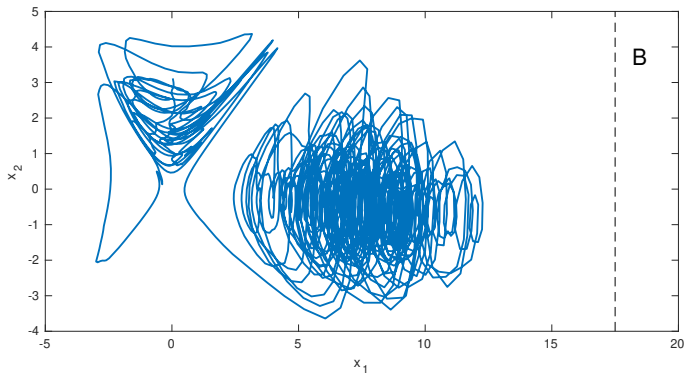
$$dy_1 = \mu \left(-L_{13}x_1 + B_{312}^1x_1x_2 + B_{311}^2x_1^2 + F_3 - \frac{\gamma_1}{\varepsilon}y_1 \right) dt + \frac{\sigma_1}{\sqrt{\varepsilon}} dW_1$$

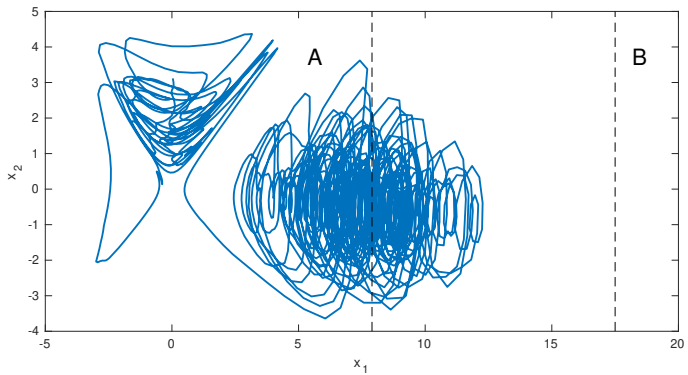
$$dy_2 = \mu \left(-L_{24}x_2 + B_{422}^3x_2x_2 + F_4 - \frac{\gamma_2}{\varepsilon}y_2 \right) dt + \frac{\sigma_2}{\sqrt{\varepsilon}} dW_2$$

$$p = \lim_{n \rightarrow \infty} \mathbb{P}(x_1 > u) = \mu(x_1 > u)$$

Franzke (2011) Model







Results

u	14	15	16	17.5	18.5
\hat{p}	$1.08 \cdot 10^{-3}$	$1.99 \cdot 10^{-4}$	$3.00 \cdot 10^{-5}$	$1.14 \cdot 10^{-6}$	$9.78 \cdot 10^{-8}$
$\text{Eff}(\hat{p})$	1.9	8.6	32.1	269.9	1521.8

Table: RMS algorithm applied to Franzke (2012) model. Parameters: $A = \{x_1 \leq 7.9\}$, $B = \{x_1 > u\}$. Importance function is $H(\mathbf{x}) = \frac{x_1}{u}$. The relative errors are below 1%.

This approach is orders of magnitude faster than Monte Carlo!

Conclusions

- ▶ We presented an algorithm for the estimation of rare events associated with the stationary distribution of a Markov chain.
- ▶ Implementation of the algorithm does not require any knowledge of the system under study – **it can be applied to ‘black-box’ models.**
- ▶ Open question: good choice of the recurrency set A .

References

1. Bisewski, K., Crommelin, D. and Mandjes, M., 2019. Rare event simulation for steady-state probabilities via recurrency cycles. *Chaos: An Interdisciplinary Journal of Nonlinear Science*, 29(3), p.033131.
2. Franzke, C., 2012. Predictability of extreme events in a nonlinear stochastic-dynamical model. *Physical Review E*, 85(3), p.031134.

Finishing remarks

Choice of the recurrency set A :

- (i) Recall that A should be such that for all cycle origins X_1^A :

$$\mathbb{P}(\tau_B < \tau_A \mid X_1^A) \equiv \mathbb{P}(\tau_B < \tau_A)$$

- (ii) At the same time A should be such that $\mathbb{E}L_1$ is not 'too large', so A should not be 'too small' ($\mu(A) \approx 0$) nor 'too big' ($\mu(A) \approx 1$).

Numerical implementation:

1. Estimate $\mathbb{E}L_1$ using Monte Carlo method. Store the locations of the cycle origins in the set $\mathcal{S} := \{X_{S_0}, X_{S_1}, \dots\}$.
2. Estimate T_B using Multilevel Splitting. Bootstrap cycle origins from the set \mathcal{S} .

Optimal Parameters

We aim to minimize the **computational time** of the algorithm under the constraint $\text{RE}^2(\hat{T}_B) = \frac{\text{Var} T_B}{(\mathbb{E} \hat{T}_B)^2} < q^2$ for a chosen $q > 0$.

$$m = c \lfloor \log p \rfloor,$$

$$p_k = \frac{2c - 1}{2c} \approx \frac{1}{5}, \quad k \in \{1, \dots, m\},$$

$$n_k = 1/p_{k+1} \approx 5, \quad k \in \{1, \dots, m - 1\},$$

$$n_m = \text{RE}(R_+) \cdot \frac{2c}{\sqrt{2c - 1}},$$

$$n_0 = \frac{1}{q\sqrt{2c - 1}} \cdot \left(\frac{c \lfloor \log p \rfloor}{\sqrt{2c - 1}} + \text{RE}(R_+) \right),$$

$$W(\hat{T}_B) \propto \frac{1}{q} \left(\frac{c \lfloor \log p \rfloor}{\sqrt{2c - 1}} + \text{RE}(R_+) \right)^2.$$

with $c \approx 0.6275$ solving $\exp(1/c) = 2c/(2c - 1)$.

How do we choose the importance function?

Let $H : \mathbb{R}^d \rightarrow [0, 1]$ be the **importance function** and for *levels* $0 = l_0 < l_1 < \dots < l_m = 1$ we put

$$B_k := \{x \in \mathbb{R}^d : H(x) \geq l_k\}.$$

Ideal importance function H should satisfy:

$$H(x) \geq H(y) \implies \mathbb{P}_x(\tau_B < \tau_A) \geq \mathbb{P}_y(\tau_B < \tau_A).$$

In particular

$$H(x) := \mathbb{P}(\tau_B < \tau_A \mid X_0 = x)$$

satisfies the above and so is

$$H_g(x) := g(H(x))$$

for any increasing function g .